

procedure derivesToLambda($L \in N$ a non-terminal of grammar,
 T a stack)
returns true if $\exists \{p_i\} \subseteq P$ permitting $L \Rightarrow^* \lambda$.

Given P the production rules of a reduced grammar and the stack T is empty on the first call of procedure.

```
foreach (  $p \in P$  with LHS  $L$  ) do (  
    if(  $p$  in  $T$  ) then (  
        continue loop  
    )  
    if (  $p$  is  $L \rightarrow \lambda$  ) then (  
        return true  
    )  
    if ( RHS of  $p$  contains a terminal ) then (  
        continue loop  
    )  
    let allderivelambda  $\leftarrow$  true  
    foreach (  $X_i \in N$  in RHS of  $p$  ) do (  
        push  $p$  onto  $T$   
        allderivelambda  $\leftarrow$  derivesToLambda( $X_i, T$ )  
        pop  $T$   
        # it takes just one  $X_i \not\Rightarrow^* \lambda$  to mean LHS of  $p \not\Rightarrow^* \lambda$   
        if( allderivelambda is false ) then (  
            break loop  
        )  
    )  
    )  
    if( allderivelambda is true ) then (  
        return true  
    )  
)  
return false
```